

AMENDMENTS TO THE CLAIMS:

This listing of claims will replace all prior versions, and listings, of claims in the application:

LISTING OF CLAIMS:

1. (Currently Amended) A cryptographic method during which an integer division of the type $q = a \text{ div } b$ and $r = a \text{ mod } b$ is performed, ~~with where q is a quotient, a is a number of containing m bits, b is a number of containing n bits, with n less than or equal to m and b_{n-1} is non-zero, b_{n-1} being the most significant bit of b, a method during which, at each iteration of a loop subscripted by i varying between 1 and $m-n+1$, comprising the following steps:~~

(i) performing a partial division of a word A, comprising of n bits of the number a, by the number b is performed in order to obtain a bit of the quotient q, wherein at least one of the numbers a and b comprises secret data;

~~the method being characterised in that~~ (ii) repeating step (i) for $m-n+1$ iterations with the same operations are being performed at each iteration, whatever regardless of the value of the quotient bit obtained, to obtain the quotient q; and

(iii) generating encrypted or decrypted data in accordance with said quotient.

2. (Currently Amended) A method according to Claim 1, ~~during which wherein~~, at each iteration, an addition of the number b to the word A and a subtraction of the number b from the word A are performed.

3. (Currently Amended) A method according to ~~one of Claims 1 to 2, during which claim 1, wherein~~ all the following steps are performed :

Input : $a = (0, a_{m-1}, \dots, a_0)$

$b = (b_{n-1}, \dots, b_0)$

Output: $q = a \text{ div } b$ and $r = a \text{ mod } b$

$A = (0, a_{m-1}, \dots, a_{m-n+1})$; $\sigma' < - 1$

For $j = 1$ to $(m-n+1)$, do:

$a < - \text{SHL}_{m+1}(a, 1)$; $\sigma < - \text{carry}$

$A < - (\sigma') \text{SUB}_n(A, b) + (\neg \sigma') \text{ADD}_n(A, b)$

$\sigma < - (\sigma' \text{ AND } \sigma') / (\sigma' \text{ AND } \text{carry}) / (\sigma' \text{ AND } \text{carry})$

$\text{lsb}(a) \ \sigma'$

$\sigma' < - \sigma$

End For

if $(\neg \sigma = \text{TRUE})$ then $A < - \text{ADD}_n(A, b)$

4. (Currently Amended) A method according to Claim 1, ~~during which wherein~~, at each iteration, ~~an operation of addition~~ either of the number b or of a number \bar{b} complementary to the number b ~~with is added to the word A is performed~~.

5. (Currently Amended) A method according to Claim 4, ~~during which further including the step~~, at each iteration, ~~an of updating is also carried out of~~ a first variable (σ') indicating whether, during the following iteration, the number b or the number \bar{b} ~~must is to~~ be added with the word A according to the quotient bit produced $(\text{lsb}(a))$.

6. (Currently Amended) A method according to Claim 4 or ~~Claim 5~~, ~~during which wherein~~ all the following steps are performed :

Input : $a = (0, a_{m-1}, \dots, a_0)$

$b = (b_{n-1}, \dots, b_0)$

Output: $q = a \text{ div } b$ and $r = a \text{ mod } b$

$A = (0, a_{m-1}, \dots, a_{m-n+1})$; $\sigma' < - 1$; $\bar{b} < - \text{CPL2}_N(b)$

For $j = 1$ to $(m-n+1)$, do:

$a < - \text{SHL}_{m+1}(a, 1)$; $\sigma < - \text{carry}$

$d_{\text{addr}} < - b_{\text{addr}} + \sigma' (\bar{b}_{\text{addr}} - b_{\text{addr}})$

$A < - \text{ADD}_n(A, d)$

$\sigma < - (\sigma' \text{ AND } \sigma') / (\sigma' \text{ AND } \text{carry}) / (\sigma' \text{ AND } \text{carry})$

lsb(a) <- σ'

σ' <- σ

End For

if ($\neg\sigma$ = TRUE) then A <- ADD_n(A, b)

7. (Currently Amended) A method according to Claim 1, during which further including the steps, at each iteration, of performing an operation of complement to 2^n of an updated data item (b or \bar{b}) or of a notional data item (c or \bar{c}) is performed, and then an operation of addition of adding the updated data item with the word A.

8. (Currently Amended) A method according to Claim 7, during which further including the step, at each iteration, an operation of updating a second variable (δ) is also performed, indicating whether, during the following iteration, the operation of complement to 2^n must is to be performed on the updated data item or on the notional data item.

9. (Currently Amended) A method according to ~~one of Claims 7 or 8, in which~~ claim 7, further including the step, at each iteration, there is also performed an of updating of a third variable (β) indicating whether the updated data item is equal to the data item b or to its complement to 2^n .

10. (Currently Amended) A method according to ~~one of Claims 7 to 9, during which~~ claim 7, wherein all the following steps are also performed :

Input : a = (0, a_{m-1}, ..., a₀)

b = (b_{n-1}, ..., b₀)

Output: q = a div b and r = a mod b

$\sigma' <- 1$; $\beta <- 1$, $\gamma <- 1$; A = (0, a_{m-1}, ..., a_{m-n+1})

for j = 1 to (m-n+1), do:

a <- SHL_{m+1}(a, 1) ; $\sigma <-$ carry

$\delta <- \sigma' / \beta$

$d_{addr} <- b_{addr} + \delta (c_{addr} - b_{addr})$

d <- CPL2_n(d)

$A <- \text{ADD}_n(A, b)$

$\sigma <- (\sigma \text{ AND } \sigma') / (\sigma \text{ AND } \text{carry}) / (\sigma' \text{ AND } \text{carry})$

$\beta <- \neg\sigma' ; \gamma <- \gamma / \delta ; \sigma' <- \sigma$

$\text{lsb}(a) = \sigma$

end for

if ($\neg\sigma = \text{TRUE}$) then $A <- \text{ADD}_n(A, b)$

11. (Currently Amended) A method according to Claim 10, ~~during which~~ wherein, at the end, the following operations are performed :

if ($\neg\beta = \text{TRUE}$) then $b <- \text{CPL2}_n(b)$

if ($\neg\gamma = \text{TRUE}$) then $c <- \text{CPL2}_n(c)$.

12. (Currently Amended) An electronic component comprising calculation means programmed to implement a method according to ~~one of Claims 1 to 11, the claim 1, said~~ calculation means comprising ~~in particular~~ a central unit associated with a memory comprising several registers for storing the data a and b.

13. (Currently Amended) A chip card comprising an ~~integrated circuit~~ electronic component according to Claim 12.